

Technical Field of the Invention

Background of the Invention

Memory devices are typically provided as internal storage areas in the computer. The term memory identifies data storage that comes in the form of integrated circuit chips. There are several different types of memory. One type is RAM (random-access memory). This is typically used as main memory in a computer environment. RAM refers to read and write memory; that is, you can both write data into RAM and read data from RAM. This is in contrast to ROM, which permits you only to read data. Most RAM is volatile, which means that it requires a steady flow of electricity to maintain its contents. As soon as the power is turned off, whatever data was in RAM is lost.

Computers almost always contain a small amount of read-only memory (ROM) that holds instructions for starting up the computer. Unlike RAM, ROM cannot be written to. An EEPROM (electrically erasable programmable read-only memory) is a special type non-volatile ROM that can be erased by exposing it to an electrical charge. Like other types of ROM, EEPROM is traditionally not as fast as RAM. EEPROM comprise a large number of memory cells having electrically isolated gates (floating gates). Data is stored in the memory cells in the form of charge on the floating gates. Charge is transported to or removed from the floating gates by programming and erase operations, respectively.

Yet another type of non-volatile memory is a Flash memory. A Flash memory is a type of EEPROM that can be erased and reprogrammed in blocks instead of one byte at a time. Many modern PCs have their BIOS stored on a flash memory chip so that it can easily be updated if necessary. Such a BIOS is sometimes called a flash BIOS. Flash

memory is also popular in modems because it enables the modem manufacturer to support new protocols as they become standardized.

A typical Flash memory comprises a memory array that includes a large number of memory cells arranged in row and column fashion. Each of the memory cells includes a floating gate field-effect transistor capable of holding a charge. The cells are usually grouped into blocks. Each of the cells within a block can be electrically programmed in a random basis by charging the floating gate. The charge can be removed from the floating gate by a block erase operation. The data in a cell is determined by the presence or absence of the charge in the floating gate.

A synchronous DRAM (SDRAM) is a type of DRAM that can run at much higher clock speeds than conventional DRAM memory. SDRAM synchronizes itself with a CPU's bus and is capable of running at 100 MHZ, about three times faster than conventional FPM (Fast Page Mode) RAM, and about twice as fast EDO (Extended Data Output) DRAM and BEDO (Burst Extended Data Output) DRAM. SDRAM's can be accessed quickly, but are volatile. Many computer systems are designed to operate using SDRAM, but would benefit from non-volatile memory.

For the reasons stated above, and for other reasons stated below which will become apparent to those skilled in the art upon reading and understanding the present specification, there is a need in the art for a non-volatile memory device that can operate in a manner similar to SDRAM operation. Further, there is a need for a memory that can perform concurrent write and read operations without contention.

Summary of the Invention

The above-mentioned problems with memory devices and other problems are addressed by the present invention and will be understood by reading and studying the following specification.

In one embodiment, a non-volatile synchronous memory device comprises an array of memory cells arranged in a plurality of addressable banks. Each bank comprises addressable rows and columns of non-volatile memory cells. A plurality of bank buffers are coupled to each of the plurality of addressable banks. Each of the plurality of bank

buffers comprises bits to store data from a row of memory cells contained in a corresponding bank of the plurality of addressable banks.

5 A method of operating a synchronous flash memory is provided in one embodiment. The method comprises storing instruction code in each of a plurality of array blocks of the synchronous flash memory, and copying the instruction code from a first array block to a buffer circuit, during a write operation, so that the instruction code can be read from the buffer circuit during the write operation.

10 In another embodiment, a method of writing to a flash memory comprises copying first data stored in a row of an array bank to a buffer circuit, performing a write operation to write second data to the array bank, and reading the first data from buffer circuit while performing the write operation.

15 In yet another embodiment, a method of operating a flash memory comprises copying first data stored in a first row of a first array bank to a buffer circuit in response to a command from an external processor coupled to the flash memory. The method further comprises performing a write operation to write second data to a second row of the first array bank in response to a write command provided by the processor, and reading the first data from buffer circuit while performing the write operation in response to a read command provided by the processor. The first data contains instruction code for the processor. A status register of the flash memory is monitored with the processor in
20 response to the instruction code.

Brief Description of the Drawings

Figure 1A is a block diagram of a synchronous flash memory of the present invention;

25 Figure 1B is an integrated circuit pin interconnect diagram of one embodiment of the present invention;

Figure 1C is an integrated circuit interconnect bump grid array diagram of one embodiment of the present invention;

Figure 2 illustrates a mode register of one embodiment of the present invention;

Figure 3 illustrates read operations having a CAS latency of one, two and three clock cycles;

Figure 4 illustrates activating a specific row in a bank of the memory of one embodiment of the present invention;

5 Figure 5 illustrates timing between an active command and a read or write command;

Figure 6 illustrates a read command;

Figure 7 illustrates timing for consecutive read bursts of one embodiment of the present invention;

10 Figure 8 illustrates random read accesses within a page of one embodiment of the present invention;

Figure 9 illustrates a read operation followed by a write operation;

Figure 10 illustrates read burst operation that are terminated using a burst terminate command according to one embodiment of the present invention;

15 Figure 11 illustrates a write command;

Figure 12 illustrates a write followed by a read operation;

Figure 13 illustrates a power-down operation of one embodiment of the present invention;

Figure 14 illustrates a clock suspend operation during a burst read;

20 Figure 15 illustrates a memory address map of one embodiment of the memory having two boot sectors;

Figure 16 is a flow chart of a self-timed write sequence according to one embodiment of the present invention;

25 Figure 17 is a flow chart of a complete write status-check sequence according to one embodiment of the present invention;

Figure 18 is a flow chart of a self-timed block erase sequence according to one embodiment of the present invention;

Figure 19 is a flow chart of a complete block erase status-check sequence according to one embodiment of the present invention;

Figure 20 is a flow chart of a block protect sequence according to one embodiment of the present invention;

Figure 21 is a flow chart of a complete block status-check sequence according to one embodiment of the present invention;

5 Figure 22 is a flow chart of a device protect sequence according to one embodiment of the present invention;

Figure 23 is a flow chart of a block unprotect sequence according to one embodiment of the present invention;

Figure 24 illustrates the timing of an initialize and load mode register operation;

10 Figure 25 illustrates the timing of a clock suspend mode operation;

Figure 26 illustrates the timing of a burst read operation;

Figure 27 illustrates the timing of alternating bank read accesses;

Figure 28 illustrates the timing of a full-page burst read operation;

Figure 29 illustrates the timing of a burst read operation using a data mask signal;

15 Figure 30 illustrates the timing of a write operation followed by a read to a different bank;

Figure 31 illustrates the timing of a write operation followed by a read to the same bank;

Figure 32 illustrates a memory system of the present invention;

20 Figure 33 illustrates a multi-processor system of the present invention; and

Figure 34 illustrates a memory device of the present invention including bank buffer circuitry.

25 **Detailed Description of the Invention**

In the following detailed description of present embodiments, reference is made to the accompanying drawings that form a part hereof, and in which is shown by way of illustration specific embodiments in which the inventions may be practiced. These embodiments are described in sufficient detail to enable those skilled in the art to practice the invention, and it is to be understood that other embodiments may be utilized and that

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logical, mechanical and electrical changes may be made without departing from the spirit and scope of the present invention. The following detailed description is, therefore, not to be taken in a limiting sense, and the scope of the present invention is defined only by the claims.

5 Referring to Figure 1A, a block diagram of one embodiment of the present invention is described. The memory device 100 includes an array of non-volatile flash memory cells 102. The array is arranged in a plurality of addressable banks. In one embodiment, the memory contains four memory banks 104, 106, 108 and 110. Each memory bank contains addressable sectors of memory cells. The data stored in the
10 memory can be accessed using externally provided location addresses received by address register 112. The addresses are decoded using row address multiplexer circuitry 114. The addresses are also decoded using bank control logic 116 and row address latch and decode circuitry 118. To access an appropriate column of the memory, column address counter and latch circuitry 120 couples the received addresses to column decode
15 circuitry 122. Circuit 124 provides input/output gating, data mask logic, read data latch circuitry and write driver circuitry. Data is input through data input registers 126 and output through data output registers 128. Command execution logic 130 is provided to control the basic operations of the memory device. A state machine 132 is also provided to control specific operations performed on the memory arrays and cells. A status
20 register 134 and an identification register 136 can also be provided to output data.

Figure 1B illustrates an interconnect pin assignment of one embodiment of the present invention. The memory package 150 has 54 interconnect pins. The pin configuration is substantially similar to available SDRAM packages. Two interconnects specific to the present invention are RP# 152 and Vccp 154. Although the present
25 invention may share interconnect labels that appear the same as SDRAM's, the function of the signals provided on the interconnects are described herein and should not be equated to SDRAM's unless set forth herein. Figure 1C illustrates one embodiment of a memory package 160 that has bump connections instead of the pin connections of Figure 1C. The present invention, therefore, is not limited to a specific package
30 configuration.

Prior to describing the operational features of the memory device, a more detailed description of the interconnect pins and their respective signals is provided.

The input clock connection is used to provide a clock signal (CLK). The clock signal can be driven by a system clock, and all synchronous flash memory input signals are sampled on the positive edge of CLK. CLK also increments an internal burst counter and controls the output registers.

The input clock enable (CKE) connection is used to activate (HIGH state) and deactivates (LOW state) the CLK signal input. Deactivating the clock input provides POWER-DOWN and STANDBY operation (where all memory banks are idle), ACTIVE POWER-DOWN (a memory row is ACTIVE in either bank) or CLOCK SUSPEND operation (burst/access in progress). CKE is synchronous except after the device enters power-down modes, where CKE becomes asynchronous until after exiting the same mode. The input buffers, including CLK, are disabled during power-down modes to provide low standby power. CKE may be tied HIGH in systems where power-down modes (other than RP# deep power-down) are not required.

The chip select (CS#) input connection provides a signal to enable (registered LOW) and disable (registered HIGH) a command decoder provided in the command execution logic. All commands are masked when CS# is registered HIGH. Further, CS# provides for external bank selection on systems with multiple banks, and CS# can be considered part of the command code; but may not be necessary.

The input command input connections for RAS#, CAS#, and WE# (along with CAS#, CS#) define a command that is to be executed by the memory, as described in detail below. The input/output mask (DQM) connections are used to provide input mask signals for write accesses and an output enable signal for read accesses. Input data is masked when DQM is sampled HIGH during a WRITE cycle. The output buffers are placed in a high impedance (High-Z) state (after a two-clock latency) when DQM is sampled HIGH during a READ cycle. DQML corresponds to data connections DQ0-DQ7 and DQMH corresponds to data connections DQ8-DQ15. DQML and DQMH are considered to be the same state when referenced as DQM.

Address inputs 133 are primarily used to provide address signals. In the illustrated embodiment the memory has 12 lines (A0-A11). Other signals can be provided on the address connections, as described below. The address inputs are sampled during an ACTIVE command (row-address A0-A11) and a READ/WRITE command (column-address A0-A7) to select one location in a respective memory bank. The address inputs are also used to provide an operating code (OpCode) during a LOAD COMMAND REGISTER operation, explained below. Address lines A0-A11 are also used to input mode settings during a LOAD MODE REGISTER operation.

An input reset/power-down (RP#) connection 140 is used for reset and power-down operations. Upon initial device power-up, a 100 μ s delay after RP# has transitioned from LOW to HIGH is required in one embodiment for internal device initialization, prior to issuing an executable command. The RP# signal clears the status register, sets the internal state machine (ISM) 132 to an array read mode, and places the device in a deep power-down mode when LOW. During power down, all input connections, including CS# 142, are "Don't Care" and all outputs are placed in a High-Z state. When the RP# signal is equal to a VHH voltage (5V), all protection modes are ignored during WRITE and ERASE. The RP# signal also allows a device protect bit to be set to 1 (protected) and allows block protect bits of a 16 bit register, at locations 0 and 15 to be set to 0 (unprotected) when brought to VHH. The protect bits are described in more detail below. RP# is held HIGH during all other modes of operation.

Bank address input connections, BA0 and BA1 define which bank an ACTIVE, READ, WRITE, or BLOCK PROTECT command is being applied. The DQ0-DQ15 connections 143 are data bus connections used for bi-directional data communication. Referring to Figure 1B, a VCCQ connection is used to provide isolated power to the DQ connections to improved noise immunity. In one embodiment, VCCQ = Vcc or 1.8V \pm 0.15V. The VSSQ connection is used to isolated ground to DQs for improved noise immunity. The VCC connection provides a power supply, such as 3V. A ground connection is provided through the Vss connection. Another optional voltage is provided on the VCCP connection 144. The VCCP connection can be tied externally to VCC, and sources current during device initialization, WRITE and ERASE operations. That is,

writing or erasing to the memory device can be performed using a VCCP voltage, while all other operations can be performed with a VCC voltage. The Vccp connection is coupled to a high voltage switch/pump circuit 145.

The following sections provide a more detailed description of the operation of the synchronous flash memory. One embodiment of the present invention is a nonvolatile, electrically sector-erasable (Flash), programmable read-only memory containing 67,108,864 bits organized as 4,194,304 words by 16 bits. Other population densities are contemplated, and the present invention is not limited to the example density. Each memory bank is organized into four independently erasable blocks (16 total). To ensure that critical firmware is protected from accidental erasure or overwrite, the memory can include sixteen 256K-word hardware and software lockable blocks. The memory's four-bank architecture supports true concurrent operations.

A read access to any bank can occur simultaneously with a background WRITE or ERASE operation to any other bank. The synchronous flash memory has a synchronous interface (all signals are registered on the positive edge of the clock signal, CLK). Read accesses to the memory can be burst oriented. That is, memory accesses start at a selected location and continue for a programmed number of locations in a programmed sequence. Read accesses begin with the registration of an ACTIVE command, followed by a READ command. The address bits registered coincident with the ACTIVE command are used to select the bank and row to be accessed. The address bits registered coincident with the READ command are used to select the starting column location and bank for the burst access.

The synchronous flash memory provides for programmable read burst lengths of 1, 2, 4 or 8 locations, or the full page, with a burst terminate option. Further, the synchronous flash memory uses an internal pipelined architecture to achieve high-speed operation.

The synchronous flash memory can operate in low-power memory systems, such as systems operating on three volts. A deep power-down mode is provided, along with a power-saving standby mode. All inputs and outputs are low voltage transistor-transistor logic (LVTTTL) compatible. The synchronous flash memory offers substantial advances

in Flash operating performance, including the ability to synchronously burst data at a high data rate with automatic column address generation and the capability to randomly change column addresses on each clock cycle during a burst access.

In general, the synchronous flash memory is configured similar to a multi-bank
5 DRAM that operates at low voltage and includes a synchronous interface. Each of the banks is organized into rows and columns. Prior to normal operation, the synchronous flash memory is initialized. The following sections provide detailed information covering device initialization, register definition, command descriptions and device operation.

10 The synchronous flash is powered up and initialized in a predefined manner. After power is applied to VCC, VCCQ and VCCP (simultaneously), and the clock signal is stable, RP# 140 is brought from a LOW state to a HIGH state. A delay, such as a 100µs delay, is needed after RP# transitions HIGH in order to complete internal device initialization. After the delay time has passed, the memory is placed in an array read
15 mode and is ready for Mode Register programming or an executable command. After initial programming of a non-volatile mode register 147 (NVMode Register), the contents are automatically loaded into a volatile Mode Register 148 during the initialization. The device will power up in a programmed state and will not require reloading of the non-volatile mode register 147 prior to issuing operational commands. This is explained in
20 greater detail below.

The Mode Register 148 is used to define the specific mode of operation of the synchronous flash memory. This definition includes the selection of a burst length, a burst type, a CAS latency, and an operating mode, as shown in Figure 2. The Mode Register is programmed via a LOAD MODE REGISTER command and retains stored
25 information until it is reprogrammed. The contents of the Mode Register may be copied into the NVMode Register 147. The NVMode Register settings automatically load the Mode Register 148 during initialization. Details on ERASE NVMODE REGISTER and WRITE NVMODE REGISTER command sequences are provided below. Those skilled in the art will recognize that an SDRAM requires that a mode register must be externally
30 loaded during each initialization operation. The present invention allows a default mode

to be stored in the NV mode register 147. The contents of the NV mode register are then copied into a volatile mode register 148 for access during memory operations.

Mode Register bits M0-M2 specify a burst length, M3 specifies a burst type (sequential or interleaved), M4-M6 specify a CAS latency, M7 and M8 specify a
5 operating mode, M9 is set to one, and M10 and M11 are reserved in this embodiment. Because WRITE bursts are not currently implemented, M9 is set to a logic one and write accesses are single location (non-burst) accesses. The Mode Register must be loaded when all banks are idle, and the controller must wait the specified time before initiating a subsequent operation.

10 Read accesses to the synchronous flash memory can be burst oriented, with the burst length being programmable, as shown in Table 1. The burst length determines the maximum number of column locations that can be automatically accessed for a given READ command. Burst lengths of 1, 2, 4, or 8 locations are available for both sequential and the interleaved burst types, and a full-page burst is available for the sequential type.
15 The full-page burst can be used in conjunction with the BURST TERMINATE command to generate arbitrary burst lengths that is, a burst can be selectively terminated to provide custom length bursts. When a READ command is issued, a block of columns equal to the burst length is effectively selected. All accesses for that burst take place within this block, meaning that the burst will wrap within the block if a boundary is reached. The
20 block is uniquely selected by A1-A7 when the burst length is set to two, by A2-A7 when the burst length is set to four, and by A3-A7 when the burst length is set to eight. The remaining (least significant) address bit(s) are used to select the starting location within the block. Full-page bursts wrap within the page if the boundary is reached.

Accesses within a given burst may be programmed to be either sequential or
25 interleaved; this is referred to as the burst type and is selected via bit M3. The ordering of accesses within a burst is determined by the burst length, the burst type and the starting column address, as shown in Table 1.

TABLE 1
BURST DEFINITION

Burst Length	Starting Column Address	Order of Accesses Within a Burst	
		Type = Sequential	Type = Interleaved
2	A0 0 1	0-1 1-0	0-1 1-0
4	A1 A0 0 0 0 1 1 0 1 1	0-1-2-3 1-2-3-0 2-3-0-1 3-0-1-2	0-1-2-3 1-0-3-2 2-3-0-1 3-2-1-0
8	A2 A1 A0 0 0 0 0 0 1 0 1 0 0 1 1 1 0 0 1 0 1 1 1 0 1 1 1	0-1-2-3-4-5-6-7 1-2-3-4-5-6-7-0 2-3-4-5-6-7-0-1 3-4-5-6-7-0-1-2 4-5-6-7-0-1-2-3 5-6-7-0-1-0-3-2 6-7-0-1-2-3-4-5 7-0-1-2-3-4-5-6	0-1-2-3-4-5-6-7 1-0-3-2-5-4-7-6 2-3-0-1-6-7-4-5 3-2-1-0-7-6-5-4 4-5-6-7-0-1-2-3 5-4-7-6-1-0-3-2 6-7-4-5-2-3-0-1 7-6-5-4-3-2-1-0
Full Page 256	n= A0-A7 (location 0-255)	Cn, Cn+1, Cn+2 Cn+3, Cn+4 ...Cn-1, Cn...	Not supported

Column Address Strobe (CAS) latency is a delay, in clock cycles, between the registration of a READ command and the availability of the first piece of output data on the DQ connections. The latency can be set to one, two or three clocks cycles. For example, if a READ command is registered at clock edge n , and the latency is m clocks, the data will be available by clock edge $n + m$. The DQ connections will start driving

data as a result of the clock edge one cycle earlier ($n + m - 1$) and, provided that the relevant access times are met, the data will be valid by clock edge $n + m$. For example, assuming that the clock cycle time is such that all relevant access times are met, if a READ command is registered at T0, and the latency is programmed to two clocks, the DQs will start driving after T1 and the data will be valid by T2, as shown in Figure 3. Figure 3 illustrates example operating frequencies at which different clock latency setting can be used. The normal operating mode is selected by setting M7 and M8 to zero, and the programmed burst length applies to READ bursts.

The following truth tables provide more detail on the operation commands of an embodiment of the memory of the present invention. An explanation is provided herein of the commands and follows Truth Table 2.

TRUTH TABLE 1

Interface Commands and DQM Operation

NAME (FUNCTION)	CS #	RAS #	CAS #	WE #	DQM	ADDR	DQs
COMMAND INHIBIT (NOP)	H	X	X	X	X	X	X
NO OPERATION (NOP)	L	H	H	H	X	X	X
ACTIVE (Select bank and activate row)	L	L	H	H	X	Bank/ Row	X
READ (Select bank, column and start READ burst)	L	H	L	H	X	Bank/ Col	X
WRITE (Select bank,	L	H	L	L	X	Bank/	Valid

column and start WRITE)						Col	
BURST TERMINATE	L	H	H	L	X	X	Active
ACTIVE TERMINATE	L	L	H	L	X	X	X
LOAD COMMAND REGISTER	L	L	L	H	X	Com Code	X
LOAD MODE REGISTER	L	L	L	L	X	Op Code	X
Write Enable/Output Enable	-	-	-	-	L	-	Active
Write Inhibit/Output High-Z	-	-	-	-	H	-	High-Z

TRUTH TABLE 2
Flash Memory Command Sequences

Operation	1 st CYCLE					2nd CYCLE					3rd CYCLE				
	C M D	A D D R	A D D R	D Q	R P #	C M D	A D D R	A D D R	D Q	R P #	C M D	A D D R	A D D R	D Q	R P #
READ DEVICE Config	L C R	9 0 H	B a n k	X	H	A C T I V E	R o w	B a n k	X	H	R E A D	C A	B a n k	X	H
READ	L	7	X	X	H	A C	X	X	X	H	R	X	X	X	H

Status Register	C R	0 H				T I V E					E A D				
CLEAR Status Register	L C R	5 0 H	X	X	H										
ERASE SETUP/Confirm	L C R	2 0 H	B a n k	X	H	A C T I V E	R o w	B a n k	X	H	W R I T E	X	B a n k	D 0 H	H / V H H
WRITE SETUP/WRITE	L C R	4 0 H	B a n k	X	H	A C T I V E	R o w	B a n k	X	H	W R I T E	C o l	B a n k	D I N	H / V H H
Protect BLOCK/Confirm	L C R	6 0 H	B a n k	X	H	A C T I V E	R o w	B a n k	X	H	W R I T E	X	B a n k	0 1 H	H / V H H
Protect DEVICE/Confirm	L C R	6 0 H	B a n k	X	H	A C T I V E	X	B a n k	X	H	W R I T E	X	B a n k	F 1 H	V H H
Unprotect BLOCKS/Confirm	L C R	6 0 H	B a n k	X	H	A C T I V E	X	B a n k	X	H	W R I T E	X	B a n k	D 0 H	H / V H H
ERASE NVmode Register	L C R	3 0 H	B a n k	X	H	A C T I V E	X	B a n k	X	H	W R I T E	X	B a n k	C 0 H	H
WRITE NVmode Register	L C R	A 0 H	B a n k	X	H	A C T I V E	X	B a n k	X	H	W R I T E	X	B a n k	X	H

The COMMAND INHIBIT function prevents new commands from being executed by the synchronous flash memory, regardless of whether the CLK signal is enabled. The synchronous flash memory is effectively deselected, but operations already in progress are not affected.

The NO OPERATION (NOP) command is used to perform a NOP to the synchronous flash memory that is selected (CS# is LOW). This prevents unwanted commands from being registered during idle or wait states, and operations already in progress are not affected.

The mode register data is loaded via inputs A0-A11. The LOAD MODE REGISTER command can only be issued when all array banks are idle, and a subsequent executable command cannot be issued until a predetermined time delay (MRD) is met. The data in the NVMode Register 147 is automatically loaded into the Mode Register 148 upon power-up initialization and is the default data unless dynamically changed with the LOAD MODE REGISTER command.

An ACTIVE command is used to open (or activate) a row in a particular array bank for a subsequent access. The value on the BA0, BA1 inputs selects the bank, and the address provided on inputs A0-A11 selects the row. This row remains active for accesses until the next ACTIVE command, power-down or RESET.

The READ command is used to initiate a burst read access to an active row. The value on the BA0, BA1 inputs selects the bank, and the address provided on inputs A0-A7 selects the starting column location. Read data appears on the DQs subject to the logic level on the data mask (DQM) input that was present two clocks earlier. If a given DQM signal was registered HIGH, the corresponding DQs will be High-Z (high impedance) two clocks later; if the DQM signal was registered LOW, the DQs will provide valid data. Thus, the DQM input can be used to mask output data during a read operation.

A WRITE command is used to initiate a single-location write access on an active row. A WRITE command must be preceded by a WRITE SETUP command. The value on the BA0, BA1 inputs selects the bank, and the address provided on inputs A0-A7

selects a column location. Input data appearing on the DQs is written to the memory array, subject to the DQM input logic level appearing coincident with the data. If a given DQM signal is registered LOW, the corresponding data will be written to memory; if the DQM signal is registered HIGH, the corresponding data inputs will be ignored, and a
5 WRITE will not be executed to that word/column location. A WRITE command with DQM HIGH is considered a NOP.

An ACTIVE TERMINATE command is not required for synchronous flash memories, but can be provided to terminate a read in a manner similar to the SDRAM PRECHARGE command. The ACTIVE TERMINATE command can be issued to
10 terminate a BURST READ in progress, and may or may not be bank specific.

A BURST TERMINATE command is used to truncate either fixed-length or full-page bursts. The most recently registered READ command prior to the BURST TERMINATE command will be truncated. BURST TERMINATE is not bank specific.

The Load Command Register operation is used to initiate flash memory control
15 commands to the Command Execution Logic (CEL) 130. The CEL receives and interprets commands to the device. These commands control the operation of the Internal State Machine 132 and the read path (i.e., memory array 102, ID Register 136 or Status Register 134).

Before any READ or WRITE commands can be issued to a bank within the
20 synchronous flash memory, a row in that bank must be "opened." This is accomplished via the ACTIVE command (defined by CS#, WE#, RAS#, CAS#), which selects both the bank and the row to be activated, see Figure 4.

After opening a row (issuing an ACTIVE command), a READ or WRITE command may be issued to that row, subject to a time period (tRCD) specification,
25 tRCD (MIN) should be divided by the clock period and rounded up to the next whole number to determine the earliest clock edge after the ACTIVE command on which a READ or WRITE command can be entered. For example, a tRCD specification of 30ns with a 90 MHz clock (11.11ns period) results in 2.7 clocks, which is rounded to 3. This is reflected in Figure 5, which covers any case where $2 < \text{tRCD (MIN)}/\text{tCK} \leq 3$. (The

same procedure is used to convert other specification limits from time units to clock cycles).

A subsequent ACTIVE command to a different row in the same bank can be issued without having to close a previous active row, provided the minimum time interval between successive ACTIVE commands to the same bank is defined by tRC.

A subsequent ACTIVE command to another bank can be issued while the first bank is being accessed, which results in a reduction of total row access overhead. The minimum time interval between successive ACTIVE commands to different banks is defined by a time period tRRD.

READ bursts are initiated with a READ command (defined by CS#, WE#, RAS#, CAS#), as shown in Figure 6. The starting column and bank addresses are provided with the READ command. During READ bursts, the valid data-out element from the starting column address will be available following the CAS latency after the READ command. Each subsequent data-out element will be valid by the next positive clock edge. Upon completion of a burst, assuming no other commands have been initiated, the DQs will go to a High-Z state. A full page burst will continue until terminated. (At the end of the page, it will wrap to column 0 and continue.) Data from any READ burst may be truncated with a subsequent READ command, and data from a fixed-length READ burst may be immediately followed by data from a subsequent READ command. In either case, a continuous flow of data can be maintained. The first data element from the new burst follows either the last element of a completed burst, or the last desired data element of a longer burst that is being truncated. The new READ command should be issued x cycles before the clock edge at which the last desired data element is valid, where x equals the CAS latency minus one. This is shown in Figure 7 for CAS latencies of one, two and three; data element $n + 3$ is either the last of a burst of four, or the last desired of a longer burst. The synchronous flash memory uses a pipelined architecture and therefore does not require the $2n$ rule associated with a prefetch architecture. A READ command can be initiated on any clock cycle following a previous READ command. Full-speed, random read accesses within a page can be performed as shown in Figure 8, or each subsequent READ may be performed to a different bank.

Data from any READ burst may be truncated with a subsequent WRITE command (WRITE commands must be preceded by WRITE SETUP), and data from a fixed-length READ burst may be immediately followed by data from a subsequent WRITE command (subject to bus turnaround limitations). The WRITE may be initiated on the clock edge immediately following the last (or last desired) data element from the READ burst, provided that I/O contention can be avoided. In a given system design, there may be the possibility that the device driving the input data would go Low-Z before the synchronous flash memory DQs go High-Z. In this case, at least a single-cycle delay should occur between the last read data and the WRITE command.

The DQM input is used to avoid I/O contention as shown in Figure 9. The DQM signal must be asserted (HIGH) at least two clocks prior to the WRITE command (DQM latency is two clocks for output buffers) to suppress data-out from the READ. Once the WRITE command is registered, the DQs will go High-Z (or remain High-Z) regardless of the state of the DQM signal. The DQM signal must be de-asserted prior to the WRITE command (DQM latency is zero clocks for input buffers) to ensure that the written data is not masked. Figure 9 shows the case where the clock frequency allows for bus contention to be avoided without adding a NOP cycle.

A fixed-length or full-page READ burst can be truncated with either ACTIVE TERMINATE (may or may not be bank specific) or BURST TERMINATE (not bank specific) commands. The ACTIVE TERMINATE or BURST TERMINATE command should be issued x cycles before the clock edge at which the last desired data element is valid, where x equals the CAS latency minus one. This is shown in Figure 10 for each possible CAS latency; data element $n + 3$ is the last desired data element of a burst of four or the last desired of a longer burst.

A single-location WRITE is initiated with a WRITE command (defined by CS#, WE#, RAS#, CAS#) as shown in Figure 11. The starting column and bank addresses are provided with the WRITE command. Once a WRITE command is registered, a READ command can be executed as defined by Truth Tables 4 and 5. An example is shown in Figure 12. During a WRITE, the valid data-in is registered coincident with the WRITE command.

Unlike SDRAM, synchronous flash does not require a PRECHARGE command to deactivate the open row in a particular bank or the open rows in all banks. The ACTIVE TERMINATE command is similar to the BURST TERMINATE command; however, ACTIVE TERMINATE may or may not be bank specific. Asserting input A10
5 HIGH during an ACTIVE TERMINATE command will terminate a BURST READ in any bank. When A10 is low during an ACTIVE TERMINATE command, BA0 and BA1 will determine which bank will undergo a terminate operation. ACTIVE TERMINATE is considered a NOP for banks not addressed by A10, BA0, BA1.

Power-down occurs if clock enable, CKE is registered LOW coincident with a
10 NOP or COMMAND INHIBIT, when no accesses are in progress. Entering power-down deactivates the input and output buffers (excluding CKE) after internal state machine operations (including WRITE operations) are completed, for power savings while in standby.

The power-down state is exited by registering a NOP or COMMAND INHIBIT
15 and CKE HIGH at the desired clock edge (meeting tCKS). See Figure 13 for an example power-down operation.

A clock suspend mode occurs when a column access/ burst is in progress and CKE is registered LOW. In the clock suspend mode, an internal clock is deactivated, "freezing" the synchronous logic. For each positive clock edge on which CKE is sampled
20 LOW, the next internal positive clock edge is suspended. Any command or data present on the input pins at the time of a suspended internal clock edge are ignored, any data present on the DQ pins will remain driven, and burst counters are not incremented, as long as the clock is suspended (see example in Figure 14). Clock suspend mode is exited by registering CKE HIGH; the internal clock and related operation will resume on the
25 subsequent positive clock edge.

The burst read/single write mode is a default mode in one embodiment. All WRITE commands result in the access of a single column location (burst of one), while READ commands access columns according to the programmed burst length and sequence. The following Truth Table 3 illustrates memory operation using the CKE
30 signal.

TRUTH TABLE 3 - CKE

CKE_{n-1}	CKE_n	CURRENT STATE	COMMAND_n	ACTION_n
L	L	POWER-DOWN CLOCK SUSPEND	X	Maintain POWER-DOWN
			X	Maintain CLOCK-SUSPEND
L	H	POWER-DOWN CLOCK SUSPEND	COMMAND INHIBIT or NOP X	Exit POWER-DOWN Exit CLOCK SUSPEND
H	L	All Banks Idle Reading or Writing	COMMAND INHIBIT or NOP VALID	POWER-DOWN Entry CLOCK SUSPEND Entry
H	H		See Truth Table 4	

5

TRUTH TABLE 4 - Current State Bank n - Command to Bank n

CURRENT STATE	CS #	RA S#	CA S#	WE #	COMMAND/ACTION
Any	H	X	X	X	COMMAND INHIBIT (NOP/continue previous operation)
	L	H	H	H	NO OPERATION (NOP/continue previous operation)
Idle	L	L	H	H	ACTIVE (Select and activate row)
	L	L	L	H	LOAD COMMAND REGISTER
	L	L	L	L	LOAD MODE REGISTER
	L	L	H	L	ACTIVE TERMINATE
Row Active	L	H	L	H	READ (Select column and start READ burst)
	L	H	L	L	WRITE (Select column and start
	L	L	H	L	WRITE)

TRUTH TABLE 5 - Current State Bank n - Command to Bank m

CURRENT STATE	CS #	RA S#	CA S#	WE #	COMMAND/ACTION
Any	H	X	X	X	COMMAND INHIBIT (NOP/continue previous operation)
	L	H	H	H	NO OPERATION (NOP/continue previous operation)
Idle	X	X	X	X	Any Command Otherwise Allowed to Bank m
Row Activating, Active, or Active Terminate	L	L	H	H	ACTIVE (Select and activate row)
	L	H	L	H	READ (Select column and start READ burst)
	L	H	L	L	WRITE (Select column and start WRITE)
	L	L	H	L	WRITE
	L	L	L	H	ACTIVE TERMINATE LOAD COMMAND REGISTER
READ	L	L	H	H	ACTIVE (Select and activate row)
	L	H	L	H	READ (Select column and start new READ burst)
	L	H	L	L	WRITE (Select column and start WRITE)
	L	L	H	L	WRITE
	L	L	L	H	ACTIVE TERMINATE LOAD COMMAND REGISTER
WRITE	L	L	H	H	ACTIVE (Select and activate row)
	L	H	L	H	READ (Select column and start READ burst)
	L	L	H	L	ACTIVE TERMINATE
	L	H	H	L	BURST TERMINATE
	L	L	L	H	LOAD COMMAND REGISTER

The synchronous flash memory incorporates a number of features to make it ideally suited for code storage and execute-in-place applications on an SDRAM bus. The memory array is segmented into individual erase blocks. Each block may be erased without affecting data stored in other blocks. These memory blocks are read, written and erased by issuing commands to the command execution logic 130 (CEL). The CEL

controls the operation of the Internal State Machine 132 (ISM), which completely controls all ERASE NVMODE REGISTER, WRITE NVMODE REGISTER, WRITE, BLOCK ERASE, BLOCK PROTECT, DEVICE PROTECT, UNPROTECT ALL BLOCKS and VERIFY operations. The ISM 132 protects each memory location from over-erasure and optimizes each memory location for maximum data retention. In addition, the ISM greatly simplifies the control necessary for writing the device in-system or in an external programmer.

The synchronous flash memory is organized into 16 independently erasable memory blocks that allow portions of the memory to be erased without affecting the rest of the memory data. Any block may be hardware-protected against inadvertent erasure or writes. A protected block requires that the RP# pin be driven to VHH (a relatively high voltage) before being modified. The 256K-word blocks at locations 0 and 15 can have additional hardware protection. Once a PROTECT BLOCK command has been executed to these blocks, an UNPROTECT ALL BLOCKS command will unlock all blocks except the blocks at locations 0 and 15, unless the RP# pin is at VHH. This provides additional security for critical code during in-system firmware updates, should an unintentional power disruption or system reset occur.

Power-up initialization, ERASE, WRITE and PROTECT timings are simplified by using an ISM to control all programming algorithms in the memory array. The ISM ensures protection against over-erasure and optimizes write margin to each cell. During WRITE operations, the ISM automatically increments and monitors WRITE attempts, verifies write margin on each memory cell and updates the ISM Status Register. When a BLOCK ERASE operation is performed, the ISM automatically overwrites the entire addressed block (eliminates over-erasure), increments and monitors ERASE attempts and sets bits in the ISM Status Register.

The 8-bit ISM Status Register 134 allows an external processor 200 to monitor the status of the ISM during WRITE, ERASE and PROTECT operations. One bit of the 8-bit Status Register (SR7) is set and cleared entirely by the ISM. This bit indicates whether the ISM is busy with an ERASE, WRITE or PROTECT task. Additional error information is set in three other bits (SR3, SR4 and SR5): write and protect block error,

erase and unprotect all blocks error, and device protection error. Status register bits SR0, SR1 and SR2 provide details on the ISM operation underway. The user can monitor whether a device-level or bank-level ISM operation (including which bank is under ISM control) is underway. These six bits (SR3-SR5) must be cleared by the host system. The status register is described in further detail below with reference to Table 2.

The CEL 130 receives and interprets commands to the device. These commands control the operation of the ISM and the read path (i.e., memory array, device configuration or status register). Commands may be issued to the CEL while the ISM is active.

To allow for maximum power conservation, the synchronous flash features a very low current, deep power-down mode. To enter this mode, the RP# pin 140 (reset/power-down) is taken to $V_{SS} \pm 0.2V$. To prevent an inadvertent RESET, RP# must be held at V_{SS} for 100ns prior to the device entering the reset mode. With RP# held at V_{SS} , the device will enter the deep power-down mode. After the device enters the deep power-down mode, a transition from LOW to HIGH on RP# will result in a device power-up initialize sequence as outlined herein. Transitioning RP# from LOW to HIGH after entering the reset mode but prior to entering deep power-down mode requires a 1 μ s delay prior to issuing an executable command. When the device enters the deep power-down mode, all buffers excluding the RP# buffer are disabled and the current draw is low, for example, a maximum of 50 μ A at 3.3V VCC. The input to RP# must remain at V_{SS} during deep power-down. Entering the RESET mode clears the Status Register 134 and sets the ISM 132 to the array read mode.

The synchronous flash memory array architecture is designed to allow sectors to be erased without disturbing the rest of the array. The array is divided into 16 addressable "blocks" that are independently erasable. By erasing blocks rather than the entire array, the total device endurance is enhanced, as is system flexibility. Only the ERASE and BLOCK PROTECT functions are block oriented. The 16 addressable blocks are equally divided into four banks 104, 106, 108 and 110 of four blocks each. The four banks have simultaneous read-while-write functionality. An ISM WRITE or ERASE operation to any bank can occur simultaneously to a READ operation to any other bank. The Status

Further, to protect a block, a three-cycle command sequence must be issued with the block address.

The synchronous flash memory can feature three different types of READs. Depending on the mode, a READ operation will produce data from the memory array, status register, or one of the device configuration registers. A READ to the device configuration register or the Status Register must be preceded by an LCR-ACTIVE cycle and burst length of data out will be defined by the mode register settings. A subsequent READ or a READ not preceded by an LCR-ACTIVE cycle will read the array. However, several differences exist and are described in the following section.

10 A READ command to any bank outputs the contents of the memory array. While a WRITE or ERASE ISM operation is taking place, a READ to any location in the bank under ISM control may output invalid data. Upon exiting a RESET operation, the device will automatically enter the array read mode.

15 Performing a READ of the Status Register 134 requires the same input sequencing as when reading the array, except that an LCR READ STATUS REGISTER (70H) cycle must precede the ACTIVE READ cycles. The burst length of the Status Register data-out is defined by the Mode Register 148. The Status Register contents are updated and latched on the next positive clock edge subject to CAS latencies. The device will automatically enter the array read mode for subsequent READs.

20 Reading any of the Device Configuration Registers 136 requires the same input sequencing as when reading the Status Register except that specific addresses must be issued. WE# must be HIGH, and DQM and CS# must be LOW. To read the manufacturer compatibility ID, addresses must be at 000000H, and to read the device ID, addresses must be at 000001H. Any of the block protect bits is read at the third address location within each erase block (xx0002H), while the device protect bit is read from
25 location 000003H.

The DQ pins are used either to input data to the array. The address pins are used either to specify an address location or to input a command to the CEL during the LOAD COMMAND REGISTER cycle. A command input issues an 8-bit command to the CEL

to control the operation mode of the device. A WRITE is used to input data to the memory array. The following section describes both types of inputs.

To perform a command input, DQM must be LOW, and CS# and WE# must be LOW. Address pins or DQ pins are used to input commands. Address pins not used for
5 input commands are "Don't Care" and must be held stable. The 8-bit command is input on DQ0-DQ7 or A0-A7 and is latched on the positive clock edge.

A WRITE to the memory array sets the desired bits to logic 0s but cannot change a given bit to a logic 1 from a logic 0. Setting any bits to a logic 1 requires that the entire block be erased. To perform a WRITE, DQM must be LOW, CS# and WE# must be
10 LOW, and VCCP must be tied to VCC. Writing to a protected block also requires that the RP# pin be brought to VHH. A0-A11 provide the address to be written, while the data to be written to the array is input on the DQ pins. The data and addresses are latched on the rising edge of the clock. A WRITE must be preceded by a WRITE SETUP command.

15 To simplify the writing of the memory blocks, the synchronous flash incorporates an ISM that controls all internal algorithms for the WRITE and ERASE cycles. An 8-bit command set is used to control the device. See Truth Tables 1 and 2 for a list of the valid commands.

The 8-bit ISM Status Register 134 (see Table 2) is polled to check for ERASE
20 NVMODE REGISTER, WRITE NVMODE REGISTER, WRITE, ERASE, BLOCK PROTECT, DEVICE PROTECT or UNPROTECT ALL BLOCKS completion or any related errors. Completion of an ISM operation can be monitored by issuing a READ STATUS REGISTER (70H) command. The contents of the Status Register will be output to DQ0-DQ7 and updated on the next positive clock edge (subject to CAS
25 latencies) for a fixed burst length as defined by the mode register settings. The ISM operation will be complete when SR7 = 1. All of the defined bits are set by the ISM, but only the ISM status bit is reset by the ISM. The erase/unprotect block, write/protect block, device protection must be cleared using a CLEAR STATUS REGISTER (50H) command. This allows the user to choose when to poll and clear the Status Register. For
30 example, a host system may perform multiple WRITE operations before checking the

Status Register instead of checking after each individual WRITE. Asserting the RP# signal or powering down the device will also clear the Status Register.

TABLE 2
STATUS REGISTER

5

STATUS BIT#	STATUS REGISTER BIT	DESCRIPTION
SR7	ISM STATUS 1 = Ready 0 = Busy	The ISMS bit displays the active status of the state machine when performing WRITE or BLOCK ERASE. The controlling logic polls this bit to determine when the erase and write status bits are valid.
SR6	RESERVED	Reserved for future use.
SR5	ERASE/UNPROTECT BLOCK STATUS 1 = BLOCK ERASE or BLOCK UNPROTECT error 0 = Successful BLOCK ERASE or UNPROTECT	ES is set to 1 after the maximum number of ERASE cycles is executed by the ISM without a successful verify. This bit is also set to 1 if a BLOCK UNPROTECT operation is unsuccessful. ES is only cleared by a CLEAR STATUS REGISTER command or by a RESET.
SR4	WRITE/PROTECT BLOCK STATUS 1 = WRITE or BLOCK PROTECT error 0 = Successful WRITE or BLOCK PROTECT	WS is set to 1 after the maximum number of WRITE cycles is executed by the ISM without a successful verify. This bit is also set to 1 if a BLOCK or DEVICE PROTECT operation is unsuccessful. WS is only cleared by a CLEAR STATUS REGISTER command or by a RESET.
SR2 SR1	BANKA1 ISM STATUS BANKA0 ISM STATUS	When SR0 = 0, the bank under ISM control can be decoded from BA0, BA1: [0,0] Bank0; [0,1] Bank1; [1,0] Bank2; [1,1] Bank3.
SR3	DEVICE PROTECT STATUS 1 = Device protected, invalid	DPS is set to 1 if an invalid WRITE, ERASE, PROTECT BLOCK,

	operation attempted 0 = Device unprotected or RP# condition met	PROTECT DEVICE or UNPROTECT ALL BLOCKS is attempted. After one of these commands is issued, the condition of RP#, the block protect bit and the device protect bit are compared to determine if the desired operation is allowed. Must be cleared by CLEAR STATUS REGISTER or by a RESET.
SR0	DEVICE/BANK ISM STATUS 1 = Device level ISM operation 0 = Bank level ISM operation	DBS is set to 1 if the ISM operation is a device-level operation. A valid READ to any bank of the array can immediately follow the registration of a device-level ISM WRITE operation. When DBS is set to 0, the ISM operation is a bank-level operation. A READ to the bank under ISM control may result in invalid data. SR2 and SR3 can be decoded to determine which bank is under ISM control.

The device ID, manufacturer compatibility ID, device protection status and block protect status can all be read by issuing a READ DEVICE CONFIGURATION (90H) command. To read the desired register, a specific address must be asserted. See Table 3
5 for more details on the various device configuration registers 136.

Table 3
DEVICE CONFIGURATION

DEVICE CONFIGURATION	ADDRESS	DATA	CONDITION
Manufacturer Compatibility	000000H	2CH	Manufacturer compatibility read
Device ID	000001H	D3H	Device ID read

Block Protect Bit	xx0002H xx0002H	DQ0 = 1 DQ0 = 0	Block protected Block unprotected
Device Protect Bit	000003H 000003H	DQ0 = 1 DQ0 = 0	Block protect modification prevented Block protect modification enabled

Commands can be issued to bring the device into different operational modes. Each mode has specific operations that can be performed while in that mode. Several modes require a sequence of commands to be written before they are reached. The following section describes the properties of each mode, and Truth Tables 1 and 2 list all command sequences required to perform the desired operation. Read-while-write functionality allows a background operation write or erase to be performed on any bank while simultaneously reading any other bank. For a write operation, the LCR-ACTIVE-WRITE command sequences in Truth Table 2 must be completed on consecutive clock cycles. However, to simplify a synchronous flash controller operation, an unlimited number of NOPs or COMMAND INHIBITs can be issued throughout the command sequence. For additional protection, these command sequences must have the same bank address for the three cycles. If the bank address changes during the LCR-ACTIVE-WRITE command sequence, or if the command sequences are not consecutive (other than NOPs and COMMAND INHIBITs, which are permitted), the write and erase status bits (SR4 and SR5) will be set and the operation prohibited.

Upon power-up and prior to issuing any operational commands to the device, the synchronous flash is initialized. After power is applied to VCC, VCCQ and VCCP (simultaneously), and the clock is stable, RP# is transitioned from LOW to HIGH. A delay (in one embodiment a 100µs delay) is required after RP# transitions HIGH in order to complete internal device initialization. The device is in the array read mode at the completion of device initialization, and an executable command can be issued to the device.

To read the device ID, manufacturer compatibility ID, device protect bit and each of the block protect bits, a READ DEVICE CONFIGURATION (90H) command is issued. While in this mode, specific addresses are issued to read the desired information. The manufacturer compatibility ID is read at 000000H; the device ID is read at 000001H.

5 The manufacturer compatibility ID and device ID are output on DQ0-DQ7. The device protect bit is read at 000003H; and each of the block protect bits is read at the third address location within each block (xx0002H). The device and block protect bits are output on DQ0.

Three consecutive commands on consecutive clock edges are needed to input data to the array (NOPs and Command Inhibits are permitted between cycles). In the first cycle, a LOAD COMMAND REGISTER command is given with WRITE SETUP (40H) on A0-A7, and the bank address is issued on BA0, BA1. The next command is ACTIVE, which activates the row address and confirms the bank address. The third cycle is WRITE, during which the starting column, the bank address, and data are issued. The

15 ISM status bit will be set on the following clock edge (subject to CAS latencies). While the ISM executes the WRITE, the ISM status bit (SR7) will be at 0. A READ operation to the bank under ISM control may produce invalid data. When the ISM status bit (SR7) is set to a logic 1, the WRITE has been completed, and the bank will be in the array read mode and ready for an executable command. Writing to hardware-protected blocks also

20 requires that the RP# pin be set to VHH prior to the third cycle (WRITE), and RP# must be held at VHH until the ISM WRITE operation is complete. The write and erase status bits (SR4 and SR5) will be set if the LCR-ACTIVE-WRITE command sequence is not completed on consecutive cycles or the bank address changes for any of the three cycles. After the ISM has initiated the WRITE, it cannot be aborted except by a RESET or by

25 powering down the part. Doing either during a WRITE may corrupt the data being written.

Executing an ERASE sequence will set all bits within a block to logic 1. The command sequence necessary to execute an ERASE is similar to that of a WRITE. To provide added security against accidental block erasure, three consecutive command

30 sequences on consecutive clock edges are required to initiate an ERASE of a block. In

the first cycle, LOAD COMMAND REGISTER is given with ERASE SETUP (20H) on A0-A7, and the bank address of the block to be erased is issued on BA0, BA1. The next command is ACTIVE, where A10, A11, BA0, BA1 provide the address of the block to be erased. The third cycle is WRITE, during which ERASE CONFRIM (D0H) is given on DQ0-DQ7 and the bank address is reissued. The ISM status bit will be set on the following clock edge (subject to CAS latencies). After ERASE CONFIRM (D0H) is issued, the ISM will start the ERASE of the addressed block. Any READ operation to the bank where the addressed block resides may output invalid data. When the ERASE operation is complete, the bank will be in the array read mode and ready for an executable command. Erasing hardware-protected blocks also requires that the RP# pin be set to VHH prior to the third cycle (WRITE), and RP# must be held at VHH until the ERASE is completed (SR7 = 1). If the LCR-ACTIVE-WRITE command sequence is not completed on consecutive cycles (NOPs and COMMAND INHIBITs are permitted between cycles) or the bank address changes for one or more of the command cycles, the write and erase status bits (SR4 and SR5) will be set and the operation is prohibited.

The contents of the Mode Register 148 may be copied into the NVMode Register 147 with a WRITE NVMODE REGISTER command. Prior to writing to the NVMode Register, an ERASE NVMODE REGISTER command sequence must be completed to set all bits in the NVMode Register to logic 1. The command sequence necessary to execute an ERASE NVMODE REGISTER and WRITE NVMODE REGISTER is similar to that of a WRITE. See Truth Table 2 for more information on the LCR-ACTIVE-WRITE commands necessary to complete ERASE NVMODE REGISTER and WRITE NVMODE REGISTER. After the WRITE cycle of the ERASE NVMODE REGISTER or WRITE NVMODE REGISTER command sequence has been registered, a READ command may be issued to the array. A new WRITE operation will not be permitted until the current ISM operation is complete and SR7 = 1.

Executing a BLOCK PROTECT sequence enables the first level of software/hardware protection for a given block. The memory includes a 16-bit register that has one bit corresponding to the 16 protectable blocks. The memory also has a register to provide a device bit used to protect the entire device from write and erase

00000-000000

are required to initiate a DEVICE PROTECT sequence. In the first cycle, LOAD COMMAND REGISTER is issued with a PROTECT SETUP (60H) on A0- A7, and a bank address is issued on BA0, BA1. The bank address is "Don't Care" but the same bank address must be used for all three cycles. The next command is ACTIVE. The third cycle is WRITE, during which a DEVICE PROTECT (F1H) command is issued on DQ0-DQ7, and RP# is brought to VHH. The ISM status bit will be set on the following clock edge (subject to CAS latencies). An executable command can be issued to the device. RP# must be held at VHH until the WRITE is completed (SR7 = 1). A new WRITE operation will not be permitted until the current ISM operation is complete.

Once the device protect bit is set, it cannot be reset to a 0. With the device protect bit set to a 1, BLOCK PROTECT or BLOCK UNPROTECT is prevented unless RP# is at VHH during either operation. The device protect bit does not affect WRITE or ERASE operations. Refer to Table 4 for more information on block and device protect operations.

Table 4
PROTECT OPERATIONS TRUTH TABLE

FUNCTION	RP#	CS #	DQM	WE #	Address	VccP	DQ0-DQ7
DEVICE UNPROTECTED							
PROTECT SETUP	H	L	H	L	60H	X	X
PROTECT BLOCK	H	L	H	L	BA	H	01H
PROTECT DEVICE	V _{HH}	L	H	L	X	X	F1H
UNPROTECT ALL BLOCKS	H/V _{HH}	L	H	L	X	H	D0H
DEVICE PROTECTED							
PROTECT SETUP	H or V _{HH}	L	H	L	60H	X	X

PROTECT BLOCK	V _{HH}	L	H	L	BA	H	01H
UNPROTECT ALL BLOCKS	V _{HH}	L	H	L	X	H	D0H

After the ISM status bit (SR7) has been set, the device/ bank (SR0), device protect (SR3), bankA0 (SR1), bankA1 (SR2), write/protect block (SR4) and erase/unprotect (SR5) status bits may be checked. If one or a combination of SR3, SR4, SR5 status bits has been set, an error has occurred during operation. The ISM cannot reset the SR3, SR4 or SR5 bits. To clear these bits, a CLEAR STATUS REGISTER (50H) command must be given. Table 5 lists the combinations of errors.

Table 5
STATUS REGISTER ERROR DECODE

STATUS BITS			ERROR DESCRIPTION
SR5	SR4	SR3	
0	0	0	No errors
0	1	0	WRITE, BLOCK PROTECT or DEVICE PROTECT error
0	1	1	Invalid BLOCK PROTECT or DEVICE PROTECT, RP# not valid (V _{HH})
0	1	1	Invalid BLOCK or DEVICE PROTECT, RP# not valid
1	0	0	ERASE or ALL BLOCK UNPROTECT error
1	0	1	Invalid ALL BLOCK UNPROTECT, RP# not valid (V _{HH})
1	1	0	Command sequencing error

The synchronous flash memory is designed and fabricated to meet advanced code and data storage requirements. To ensure this level of reliability, VCCP must be tied to Vcc during WRITE or ERASE cycles. Operation outside these limits may reduce the

number of WRITE and ERASE cycles that can be performed on the device. Each block is designed and processed for a minimum of 100,000-WRITE/ERASE-cycle endurance.

The synchronous flash memory offers several power-saving features that may be utilized in the array read mode to conserve power. A deep power-down mode is enabled by bringing RP# to VSS $\pm 0.2V$. Current draw (ICC) in this mode is low, such as a maximum of 50 μA . When CS# is HIGH, the device will enter the active standby mode. In this mode the current is also low, such as a maximum ICC current of 30mA. If CS# is brought HIGH during a write, erase, or protect operation, the ISM will continue the WRITE operation, and the device will consume active Iccp power until the operation is completed.

Referring to Figure 16, a flow chart of a self-timed write sequence according to one embodiment of the present invention is described. The sequence includes loading the command register (code 40H), receiving an active command and a row address, and receiving a write command and a column address. The sequence then provides for a status register polling to determine if the write is complete. The polling monitors status register bit 7 (SR7) to determine if it is set to a 1. An optional status check can be included. When the write is completed, the array is placed in the array read mode.

Referring to Figure 17, a flow chart of a complete write status-check sequence according to one embodiment of the present invention is provided. The sequence looks for status register bit 4 (SR4) to determine if it is set to a 0. If SR4 is a 1, there was an error in the write operation. The sequence also looks for status register bit 3 (SR3) to determine if it is set to a 0. If SR3 is a 1, there was an invalid write error during the write operation.

Referring to Figure 18, a flow chart of a self-timed block erase sequence according to one embodiment of the present invention is provided. The sequence includes loading the command register (code 20H), and receiving an active command and a row address. The memory then determines if the block is protected. If it is not protected, the memory performs a write operation (D0H) to the block and monitors the status register for completion. An optional status check can be performed and the memory is placed in an

array read mode. If the block is protected, the erase is not allowed unless the RP# signal is at an elevated voltage (VHH).

Figure 19 illustrates a flow chart of a complete block erase status-check sequence according to one embodiment of the present invention. The sequence monitors the status register to determine if a command sequence error occurred (SR4 or SR5 = 1). If SR3 is set to a 1, an invalid erase or unprotect error occurred. Finally, a block erase or unprotect error happened if SR5 is set to a 1.

Figure 20 is a flow chart of a block protect sequence according to one embodiment of the present invention. The sequence includes loading the command register (code 60H), and receiving an active command and a row address. The memory then determines if the block is protected. If it is not protected, the memory performs a write operation (01H) to the block and monitors the status register for completion. An optional status check can be performed and the memory is placed in an array read mode. If the block is protected, the erase is not allowed unless the RP# signal is at an elevated voltage (VHH).

Referring to Figure 21, a flow chart of a complete block status-check sequence according to one embodiment of the present invention is provided. The sequence monitors the status register bits 3, 4 and 5 to determine of errors were detected.

Figure 22 is a flow chart of a device protect sequence according to one embodiment of the present invention. The sequence includes loading the command register (code 60H), and receiving an active command and a row address. The memory then determines if RP# is at VHH. The memory performs a write operation (F1H) and monitors the status register for completion. An optional status check can be performed and the memory is placed in an array read mode.

Figure 23 is a flow chart of a block unprotect sequence according to one embodiment of the present invention. The sequence includes loading the command register (code 60H), and receiving an active command and a row address. The memory then determines if the memory device is protected. If it is not protected, the memory determines if the boot locations (blocks 0 and 15) are protected. If none of the blocks are protected the memory performs a write operation (D0H) to the block and monitors the status register for completion. An optional status check can be performed and the

memory is placed in an array read mode. If the device is protected, the erase is not allowed unless the RP# signal is at an elevated voltage (VHH). Likewise, if the boot locations are protected, the memory determines if all blocks should be unprotected.

Figure 24 illustrates the timing of an initialize and load mode register operation.

5 The mode register is programmed by providing a load mode register command and providing operation code (opcode) on the address lines. The opcode is loaded into the mode register. As explained above, the contents of the non-volatile mode register are automatically loaded into the mode register upon power-up and the load mode register operation may not be needed.

10 Figure 25 illustrates the timing of a clock suspend mode operation, and Figure 26 illustrates the timing of another burst read operation. Figure 27 illustrates the timing of alternating bank read accesses. Here active command are needed to change bank addresses. A full page burst read operation is illustrated in Figure 28. Note that the full page burst does not self terminate, but requires a terminate command.

15 Figure 29 illustrates the timing of a read operation using a data mask signal. The DQM signal is used to mask the data output so that Dout m+1 is not provided on the DQ connections.

Referring to Figure 30, the timing of a write operation followed by a read to a different bank is illustrated. In this operation, a write is performed to bank a and a
20 subsequent read is performed to bank b. The same row is accessed in each bank.

Referring to Figure 31, the timing of a write operation followed by a read to the same bank is illustrated. In this operation, a write is performed to bank a and a subsequent read is performed to bank a. A different row is accessed for the read operation, and the memory must wait for the prior write operation to be completed. This
25 is different from the read of Figure 30 where the read was not delayed due to the write operation.

The synchronous Flash memory provides for a latency free write operation. This is different from a SDRAM that requires the system to provide latency for write operations, just like a read operation. So the write operation does not take away from the system bus
30 as many cycles as the SDRAM takes, and hence can improve the system read throughput,

see Figure 12 where the write data, Din, is provided on the same clock cycle as the write command and column address. The clock cycle, T1, of Figure 12 does not need to be a NOP command (see Figure 30). The read command can be provided on the next clock cycle following the write data. Thus, while the read operation requires that the DQ connections remain available for a predetermined number of clock cycles following the read command (latency), the DQ connections can be used immediately after the write command is provided (no latency). As such, the present invention allows for zero bus turn around capability. This is substantially different from the SDRAM, where multiple waits are required on the system bus when going between read and write operations. The synchronous flash provides these two features, and could improve bus throughput.

Referring to Figure 32, a system 300 of the present invention includes a synchronous memory 302 that has internal write latches 304 that are used to store write data received on the DQ inputs 306. The write latches are coupled to the memory array 310. Again, the memory array can be arranged in a number of addressable blocks. Data can be written to one block while a read operation can be performed on the other blocks. The memory cells of the array can be non-volatile memory cells. Data communication connections 306 are used for bi-directional data communication with an external device, such as a processor 320 or other memory controller.

A data buffer 330 can be coupled to the data communication connections to manage the bi-directional data communication. This buffer can be a traditional FIFO or pipelined input/output buffer circuit. The write latch is coupled between the data buffer and the memory array to latch data provided on the data communication connections. Finally, a control circuit is provided to manage the read and write operations performed on the array.

By latching the input write data, the data bus 306 (DQ's) can be released and the write operation performed using the latched data. Subsequent write operations to the memory can be prohibited while the first write operation is being performed. The bus, however, is available to immediately perform a read operation on the memory.

The present invention should not be confused with traditional input/output buffer architecture. That is, while prior memory devices used an input buffer on the DQ input

path and an output buffer on the DQ output path, the clock latency used for both read and write operations are maintained the same. The present invention can include input/output buffer circuitry to provide an interface with the DQ connections and an external processor. The additional write latches allow the memory to isolate the write path/operation to one area of the memory while allowing data read operations on other memory areas.

Prior flash memory devices have had very limited concurrent operation capabilities. That is, prior flash memories typically blocked reading from the memory while a write operation was performed. Some memory devices allowed reading while writing by suspending an ongoing write operation, then allowing the reading of the array. Yet other flash memories allow for read-while-write by providing a limited sector that can be written while the remainder of the memory is available for reading. The intention of such flash memories has been to eliminate the need for a separate EEPROM in a system. The limited sector space provides an EEPROM location in the flash and leaves the rest of memory for the flash operation.

The present invention provides a flash array that is arranged into bank structures similar to an SDRAM. In one embodiment, a 64M synchronous flash is divided into four banks with the same addressing as a 64M SDRAM. Those banks are further broken into smaller addressable sectors that can be erased or programmed. The memory allows concurrent reading and writing on a per bank basis. Thus, one bank can be written while a concurrent read can be executed on any of the other banks.

As known to one skilled in the art, an SDRAM can have one common row open in each bank. Read and write operations can be performed sequentially over the open row and across the array banks.

The present synchronous flash memory has a bank architecture similar to SDRAM that allows a row to be opened in each of the banks while one of the banks is being written to. Referring to Figure 33, one embodiment of a processing system 400 of the present invention is illustrated. A synchronous flash memory 410 is coupled to multiple processors 440, 442, 444 and 446 via a bi-directional data bus 435. The memory includes an array of non-volatile memory cells arranged in multiple banks 412, 414, 416 and 418.

Read/write circuitry 430 is generally illustrated to manage the data communication with the array. In operation, a processor, such as processor 440, can initiate a write operation on row 420 of array bank 412. While the write operation is being performed, a second processor, such as processor 442, can read data from row 420 of a second array bank.

- 5 This allows four processors to be independently working on the synchronous flash memory device. The present invention is not limited to four banks or four processors. Single or multiple processors can perform a write operation on one bank while reading data from the remaining memory array banks, see for example Figure 32. Thus, there can be concurrent write and multiple concurrent read operations ongoing.

10

Open page Architecture

- Flash memories can be used to store code and data in a nonvolatile manner. These memories are not as often written as their counter part memories, such as DRAM and SRAM memories. The reason is that the speed of programming and erasing flash
- 15 memories tends to be slow in comparison. A typical time to program a byte of data in a flash memory is in 10's of microseconds, and an erase operation can take seconds. As such, a wait time in the order of seconds is too long for a processor to wait for important data. Thus, the present invention allows simultaneous read and write operations so that a segment of the memory is accessible for read operations while other sectors are being
- 20 written.

- In prior flash memory systems, while a bank is written to, an external processor reads a status of the memory to see if the write operation has been completed. Typically, the processor performs a loop read operation to monitor the memory status until the memory signals completion of the operation. To perform the loop operation, a code that
- 25 instructs the processor to do the loop is stored in the flash memory itself. The loop instruction is usually stored in each memory array bank. Thus, the processor can read the instruction from a second bank. One drawback with the prior flash memory approach is that a processor accessing the second bank may not be the processor assigned to the first bank. That is, multiple processors can be used to access different memory banks. It is
- 30 undesirable, therefore, to have one processor reading a second memory array bank. This

is particularly relevant with the present invention that allows concurrent read and write operations to different memory array banks.

The present invention provides a method and circuitry to allow a processor to write to one memory bank and also read processing instructions from that bank. The present invention allows a first processor to open one page (row) in a bank in response to a write operation and to read that page while the bank is written to. This allows a loop routine to be stored in one page in every bank and the processor can activate that page and open it before programming or erasing of that bank.

As explained above, one embodiment of the current synchronous flash memory has four memory banks. If one bank is being erased or programmed, the other three banks are available for reading. Thus, the flash memory can accommodate multiple processors. The present flash memory also includes a set of buffers coupled to each bank of memory. The buffers are capable of storing a full page of data from the bank of memory cells. During a read operation, data can be transferred from the flash cells to the buffer. This buffer can be easily read, and is used to provide random reads across the page. Further, the present flash memory allows the data stored in the buffer to be read during a write operation. That is, data in a buffer associated with one bank can be read while a write operation is performed on the same bank.

Prior to writing data to a memory bank, a row of data from the bank is loaded in the buffer. In addition, the data can be from any page of the memory, including the page being written to. See Figure 34 for an illustration of a system 500 having array blocks 412, 414, 416 and 418 and corresponding buffer circuits 512, 514, 516 and 518. The memory thereby allows instruction code to be stored in a row 520 of memory bank 412 and transferred to buffer 512 when a write operation is performed on the bank. External processor 440 can then read the instruction code from the buffer and perform a write monitor operation based on the buffer data. In operation of one embodiment, external processor 440 provides a command(s) to copy row 520 from first array bank 412 to an associated buffer 512. This command can be an active command accompanied by the page address.

